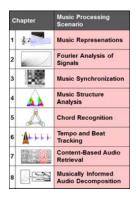


Book: Fundamentals of Music Processing



Meinard Müller

Fundamentals of Music Processing Audio, Analysis, Algorithms, Applications 483 p., 249 illus., hardcover ISBN: 978-3-319-21944-8 Springer, 2015

Accompanying website: www.music-processing.de

Chapter 3: Music Synchronization

3.1 Audio Features

- 3.2 Dynamic Time Warping
- 3.3 Applications3.4 Further Notes



As a first music processing task, we study in Chapter 3 the problem of music synchronization. The objective is to temporally align compatible representations of the same piece of music. Considering this scenario, we explain the need for musically informed audio features. In particular, we introduce the concept of chroma-based music features, which capture properties that are related to harmony and melody. Furthermore, we study an alignment technique known as dynamic time warping (DTW), a concept that is applicable for the analysis of general time series. For its efficient computation, we discuss an algorithm based on dynamic programming—a widely used method for solving a complex problem by breaking it down into a collection of simpler subproblems.

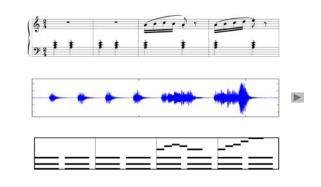
Book: Fundamentals of Music Processing

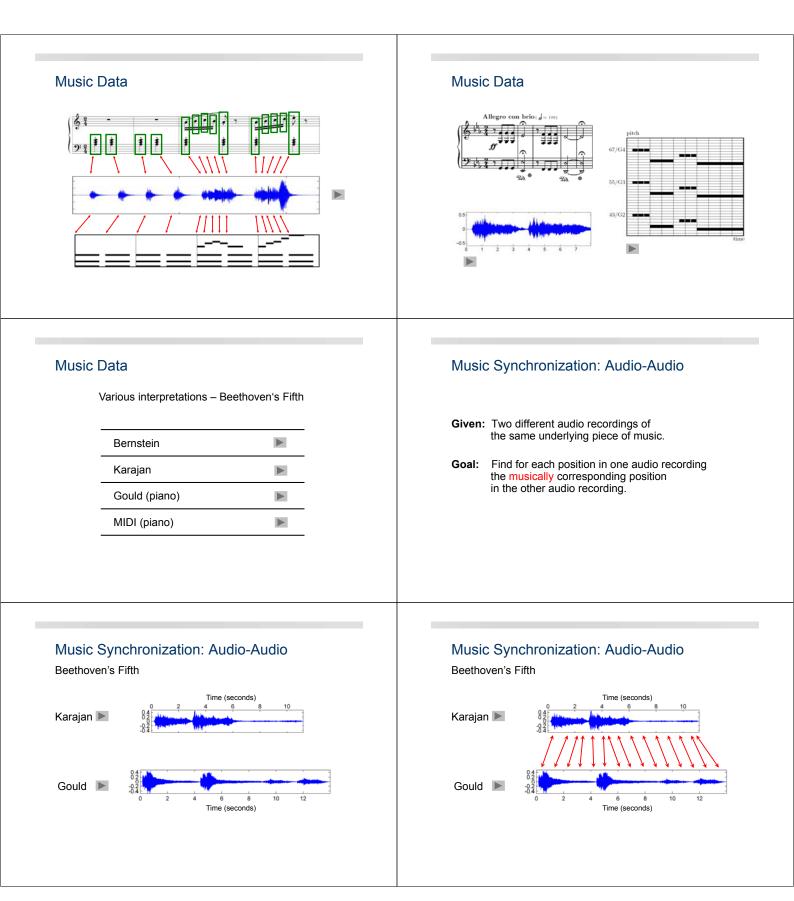
Chapter	Music Processing Scenario
1 6.2	Music Represenations
2	Fourier Analysis of Signals
3	Music Synchronization
4	Music Structure Analysis
5	Chord Recognition
6	Tempo and Beat Tracking
7	Content-Based Audio Retrieval
8	Musically Informed Audio Decomposition

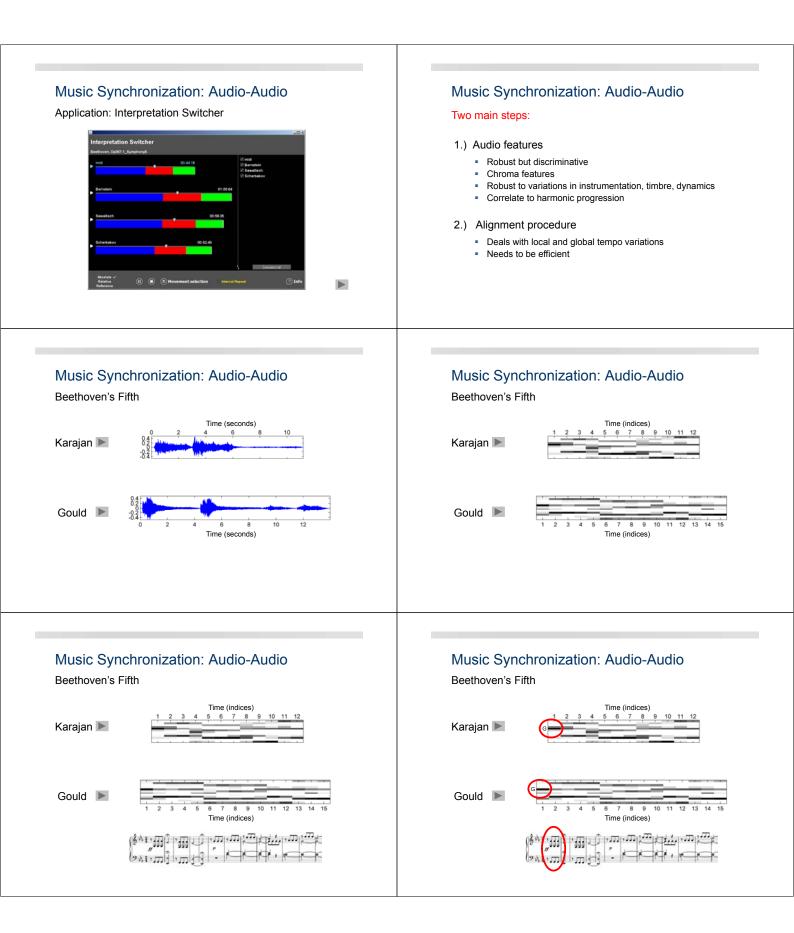
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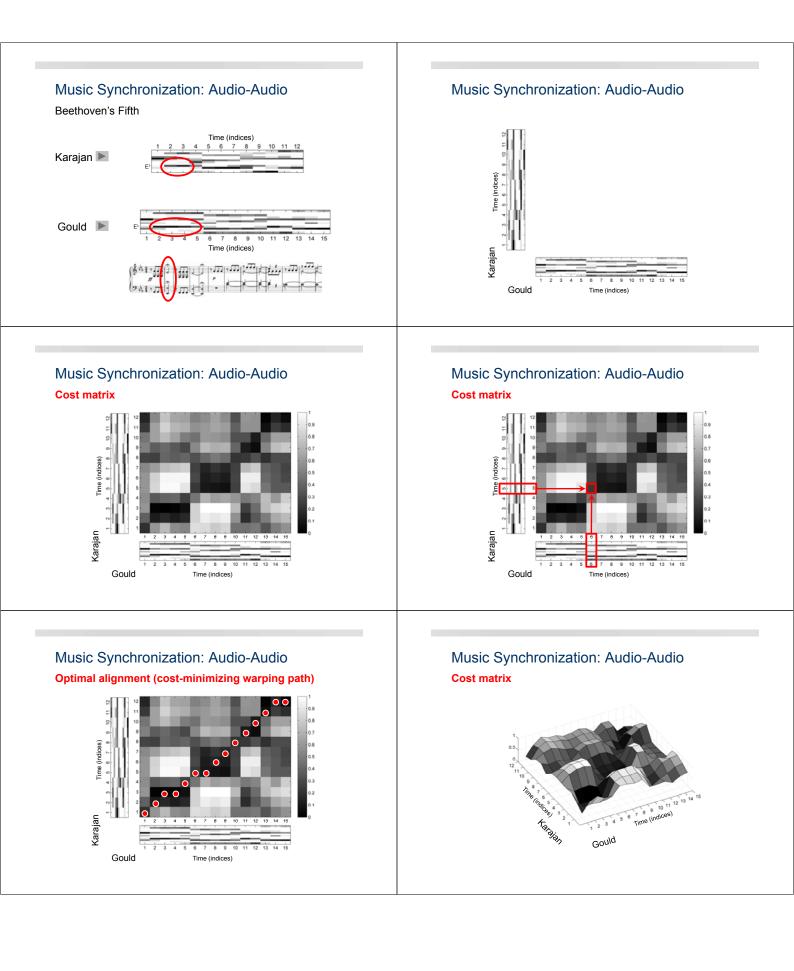
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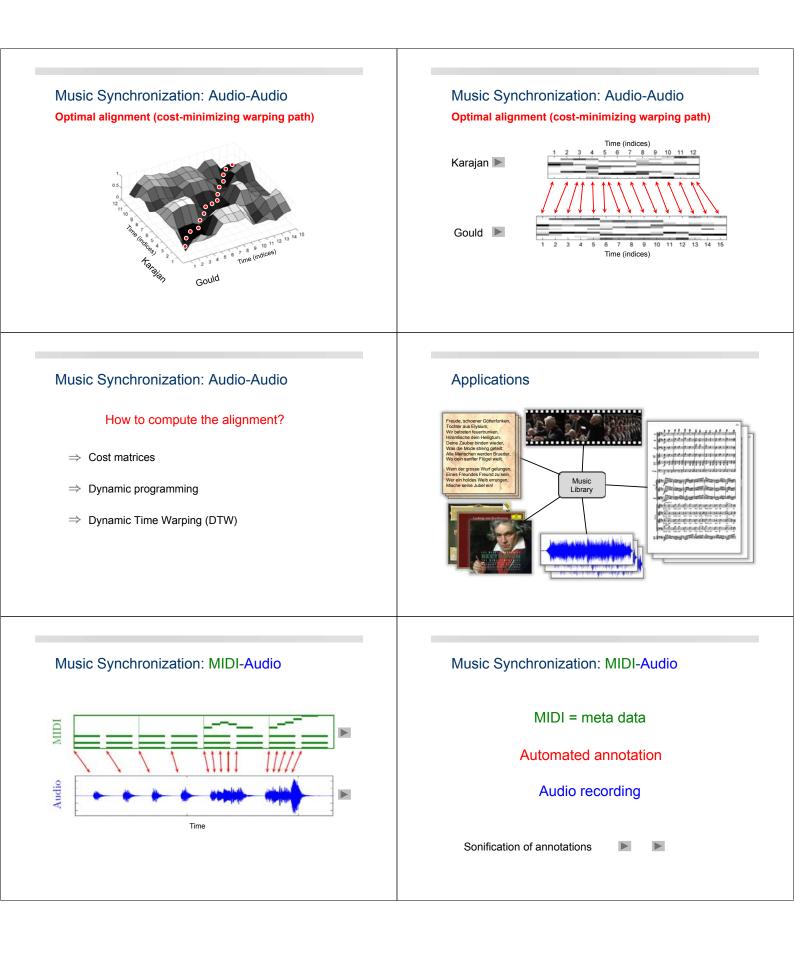
Music Data

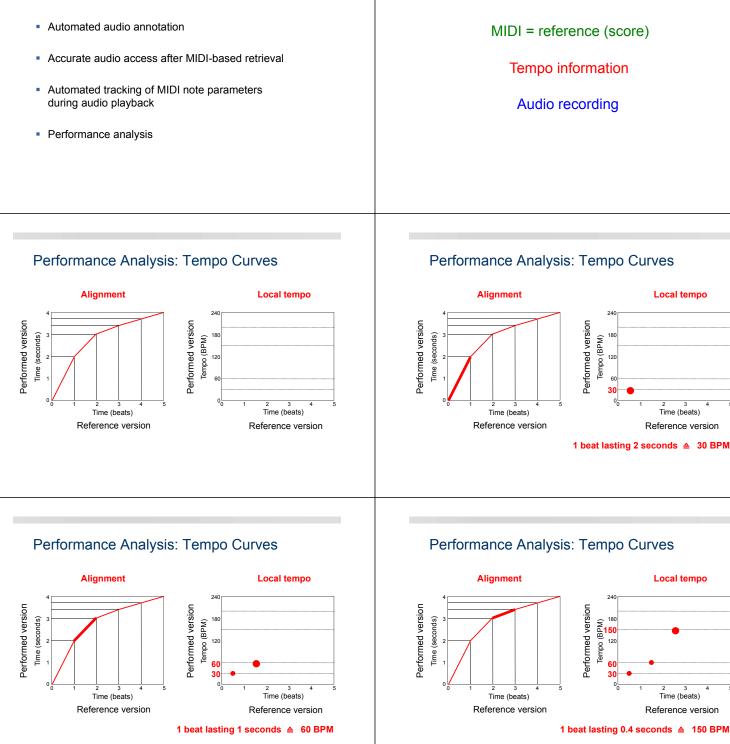






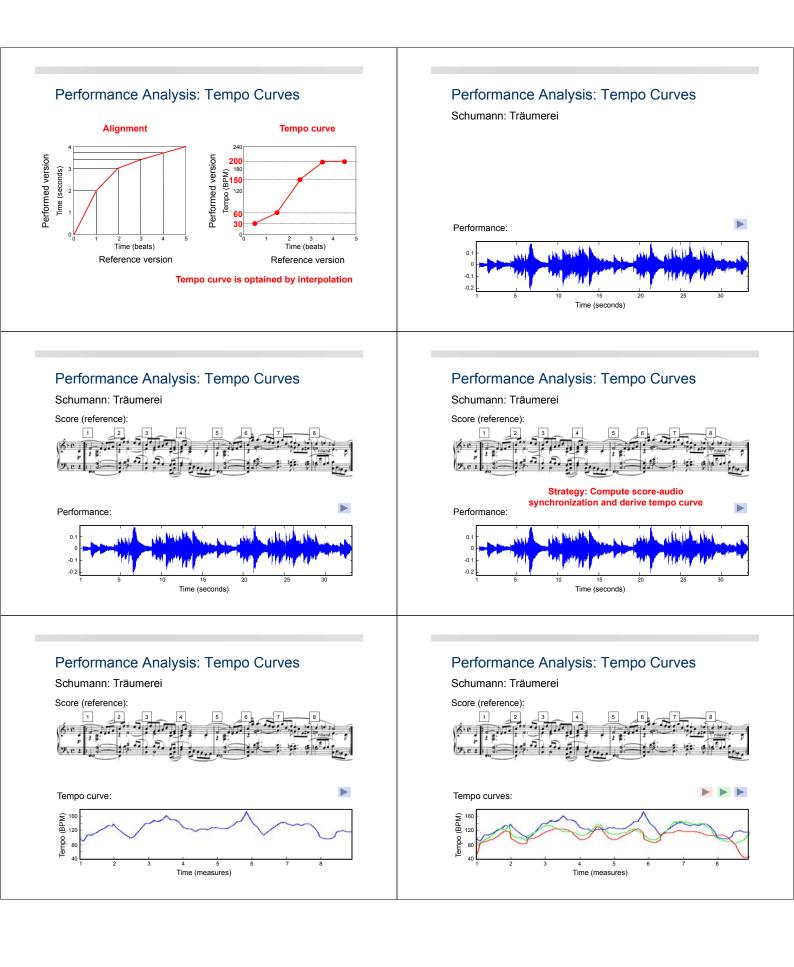


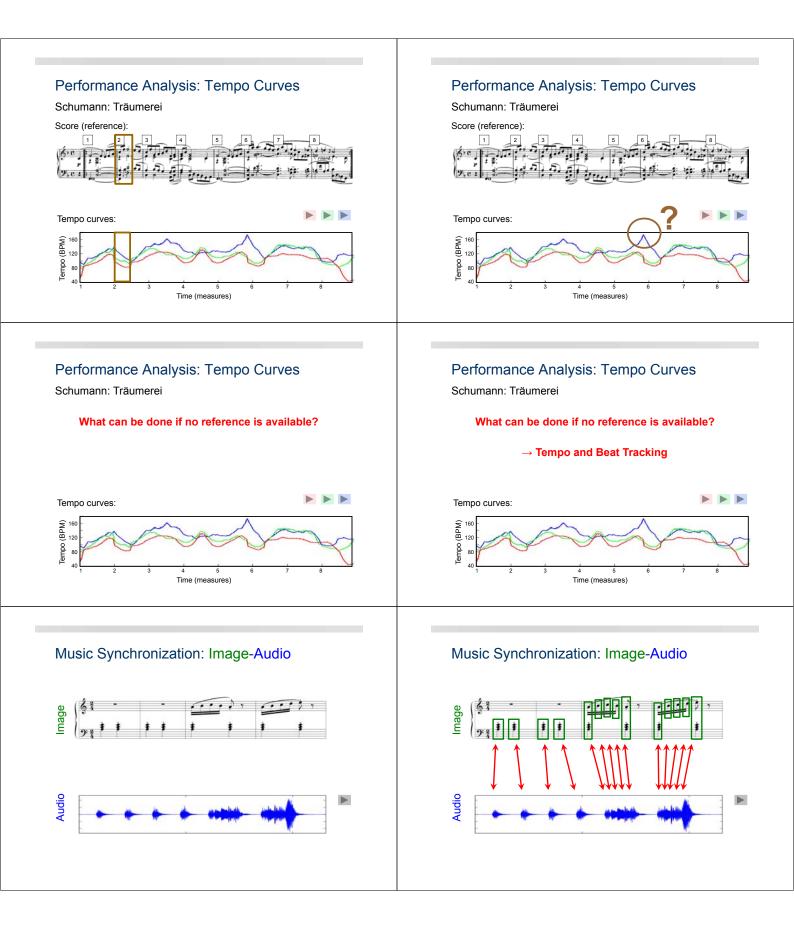


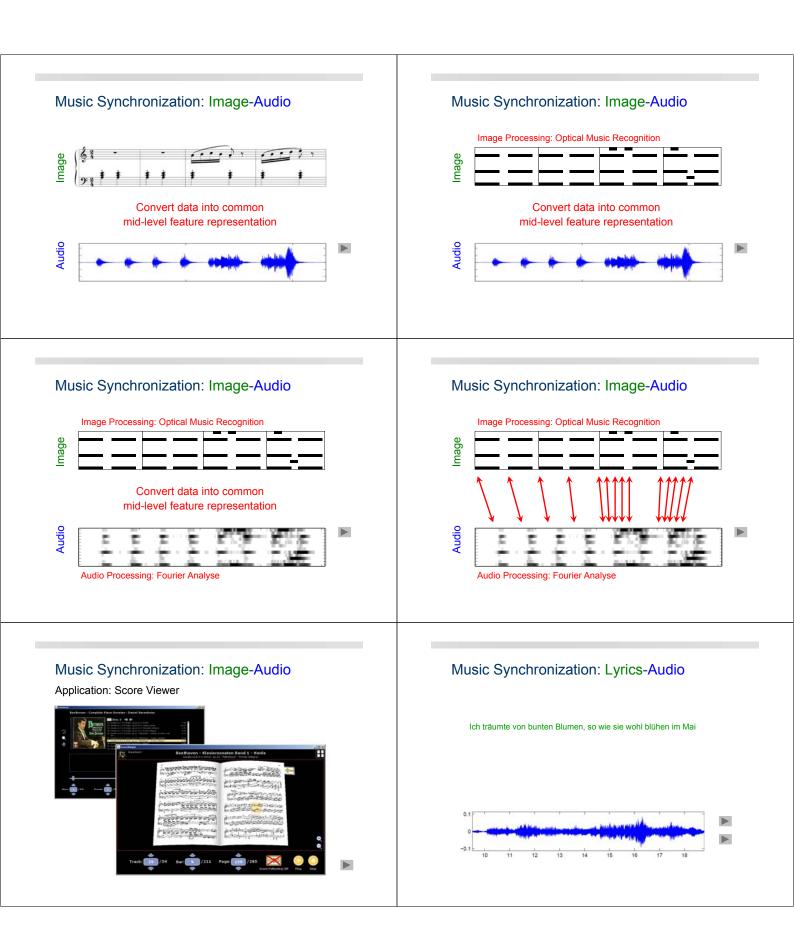


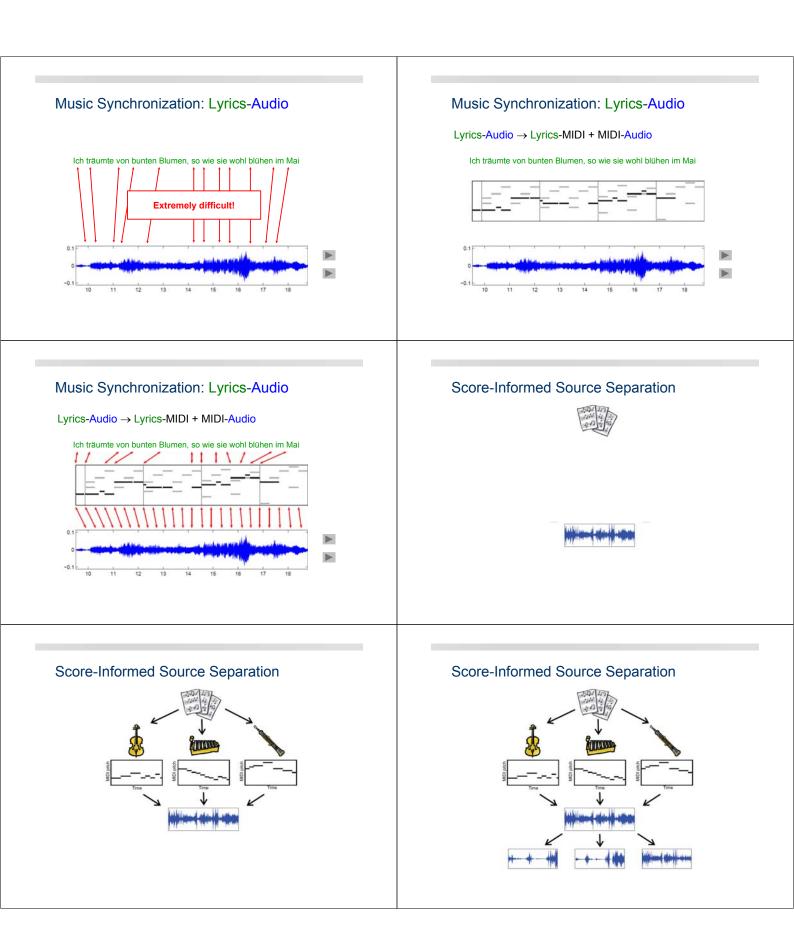
Music Synchronization: MIDI-Audio

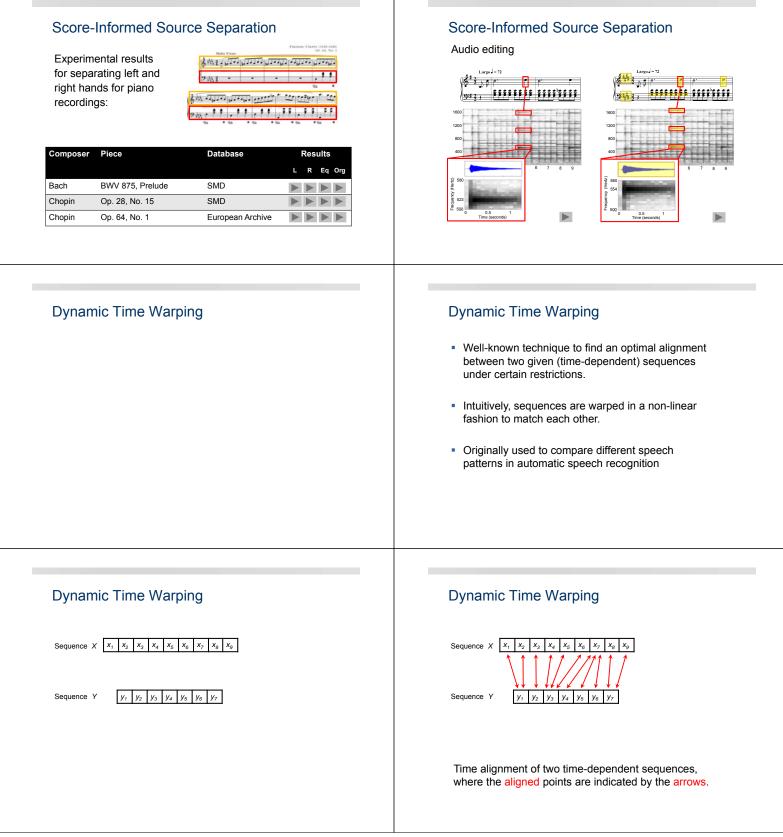
Music Synchronization: MIDI-Audio

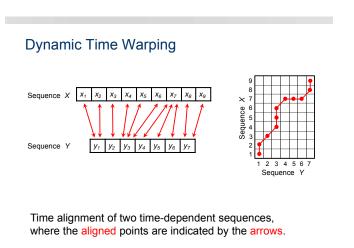












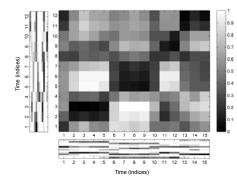
Dynamic Time Warping

To compare two different features $\ x,y\in \mathcal{F}$ one needs a local cost measure which is defined to be a function

 $c: \mathcal{F} \times \mathcal{F} \to \mathbb{R}_{\geq 0}$

Typically, c(x,y) is small (low cost) if x and y are similar to each other, and otherwise c(x,y) is large (high cost).

Dynamic Time Warping Cost matrix C



Dynamic Time Warping

The objective of DTW is to compare two (time-dependent) sequences

$$X := (x_1, x_2, \ldots, x_N)$$

of length $N \in \mathbb{N}$ and

$$Y := (y_1, y_2, \dots, y_M)$$

of length $M \in \mathbb{N}$. Here,

 $x_n, y_m \in \mathcal{F}, n \in [1:N], m \in [1:M],$

are suitable features that are elements from a given feature space denoted by $\ensuremath{\mathcal{F}}$.

Dynamic Time Warping

Evaluating the local cost measure for each pair of elements of the sequences X and Y one obtains the cost matrix

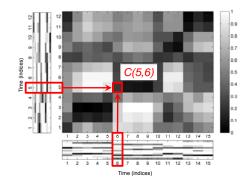
$$C \in \mathbb{R}^{N \times M}$$

denfined by

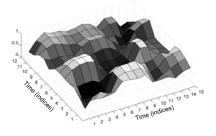
$$C(n,m) := c(x_n, y_m).$$

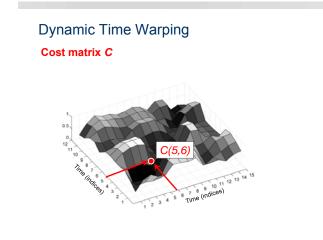
Then the goal is to find an alignment between X and Y having minimal overall cost. Intuitively, such an optimal alignment runs along a "valley" of low cost within the cost matrix C.

Dynamic Time Warping Cost matrix C



Dynamic Time Warping Cost matrix C





Dynamic Time Warping

The next definition formalizes the notion of an alignment.

A warping path is a sequence $p = (p_1, \ldots, p_L)$ with $p_{\ell} = (n_{\ell}, m_{\ell}) \in [1:N] \times [1:M]$

for $\ell \in [1:L]$ satisfying the following three conditions:

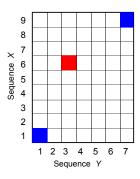
- $p_1 = (1,1)$ and $p_L = (N,M)$ Boundary condition:

• Monotonicity condition: $n_1 \leq n_2 \leq \ldots \leq n_L$ and $m_1 \leq m_2 \leq \ldots \leq m_L$

Step size condition:

 $p_{\ell+1} - p_{\ell} \in \{(1,0), (0,1), (1,1)\}$ for $\ell \in [1:L-1]$

Dynamic Time Warping Warping path

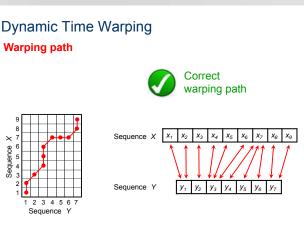


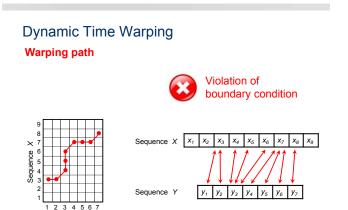
Sequence

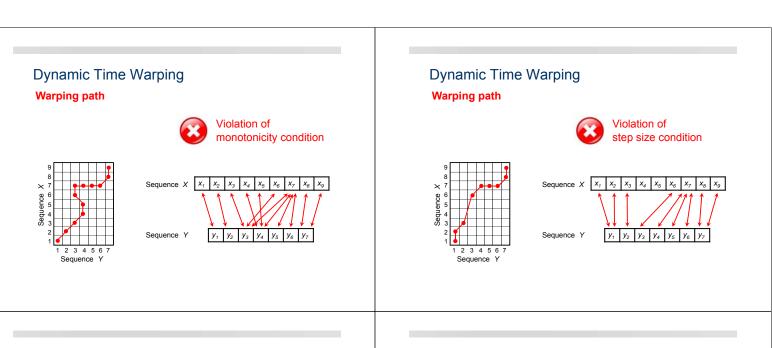
Each matrix entry (cell) corresponds to a pair of indices.

Cell = (6,3)

Boundary cells: $p_1 = (1,1)$ $p_L = (N,M) = (9,7)$







Dynamic Time Warping

The total cost $c_p(X, Y)$ of a warping path p between X and Y with respect to the local cost measure c is defined as

$$c_p(X,Y) := \sum_{\ell=1}^{n} c(x_{n_\ell}, y_{m_\ell})$$

Furthermore, an optimal warping path between X and Y is a warping path p^* having minimal total cost among all possible warping paths. The DTW distance DTW(X, Y) between X and Y is then defined as the total cost o p^*

 $DTW(X,Y) := c_{p^*}(X,Y)$ = min{c_p(X,Y) | p is a warping path}

Dynamic Time Warping

 $\begin{array}{rcl} \text{Notation:} & X(1:n) & := & (x_1, \dots, x_n), & 1 \le n \le N \\ & Y(1:m) & := & (y_1, \dots, y_m), & 1 \le m \le M \\ & D(n,m) & := & \text{DTW}(X(1:n), Y(1:m)) \end{array}$

The matrix D is called the accumulated cost matrix.

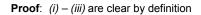
The entry D(n,m) specifies the cost of an optimal warping path that aligns X(1:n) with Y(1:m).

Dynamic Time Warping

- The warping path p^* is not unique (in general).
- DTW does (in general) not definne a metric since it may not satisfy the triangle inequality.
- There exist exponentially many warping paths.
- How can p^* be computed efficiently?

Dynamic Time Warping

Lemma:



Dynamic Time Warping

Proof of *(iv)*: Induction via n, m:

Let n > 1, m > 1 and $q = (q_1, \ldots, p_{L-1}, p_L)$ be an optimal warping path for X(1:n) and Y(1:m). Then $q_L = (n, m)$ (boundary condition).

Let $q_{L-1} = (a, b)$. The step size condition implies

 $(a,b) \in \{(n-1,m-1), (n-1,m), (n,m-1)\}$

The warping path (q_1, \ldots, q_{L-1}) must be optimal for X(1:a), Y(1:b). Thus,

 $D(n,m) = c_{(q_1,\dots,q_{L-1})}(X(1:a),Y(1:b)) + C(n,m)$

Dynamic Time Warping

Optimal warping path

Given to the algorithm is the accumulated cost matrix D. The optimal path $p^*=(p_1,\ldots,p_L)$ is computed in reverse order of the indices starting with $p_L=(N,M)$. Suppose $p_\ell=(n,m)$ has been computed. In case (n,m)=(1,1), one must have $\ell=1$ and we are done. Otherwise,

 $p_{\ell-1} := \begin{cases} (1, m-1), & \text{if } n=1\\ (n-1, 1), & \text{if } m=1\\ \arg\min\{D(n-1, m-1), \\ D(n-1, m), D(n, m-1)\}, & \text{otherwise}, \end{cases}$

where we take the lexicographically smallest pair in case "argmin" is not unique.

Dynamic Time Warping

Summary

Input:	Cost matrix C of size $N \times M$		
	Accumulated cost matrix D		
601334 * (2003	Optimal warping path P*		
Procedu	re: Initialize $(N \times$	M) matrix D by $\mathbf{D}(n,1) = \sum_{k=1}^{n} \mathbf{C}(k,1)$ for $n \in [1:N]$ and	
	$= \sum_{k=1}^{m} \mathbf{C}(1,k)$ for n	$n \in [1: M]$. Then compute in a nested loop for $n = 2,, N$ and	
		1992 19 19 19 19 19 19 19 19 19 19 19 19 19	
	$\mathbf{D}(n,m) = \mathbf{C}(n,m)$	$+\min \{ \mathbf{D}(n-1,m-1), \mathbf{D}(n-1,m), \mathbf{D}(n,m-1) \}.$	
Set $\ell = 1$	and $q_{\ell} = (N, M)$. T	hen repeat the following steps until $q_{\ell} = (1, 1)$:	
		ad let $(n,m) = q_{\ell-1}$.	
	If $n = 1$, then		
	else if $m = 1$, then		
	else	$q_i = \operatorname{argmin} \{ \mathbf{D}(n-1,m-1), \mathbf{D}(n-1,m), \mathbf{D}(n,m-1) \}.$	
		(If 'argmin' is not unique, take lexicographically smallest cell.)	
		$(L, q_{L-1}, \ldots, q_1)$ as well as D .	

Dynamic Time Warping

Accumulated cost matrix

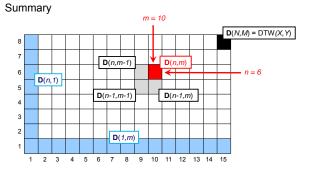
Given the two feature sequences \boldsymbol{X} and $\boldsymbol{Y}, \mbox{ the matrix } \boldsymbol{D}$ is computed recursively.

- Initialize Dusing (ii) and (iii) of the lemma.
- Compute D(n,m) for n > 1, m > 1 using (iv).
- DTW(X, Y) = D(N, M) using (i).

Note:

- Complexity O(NM).
- Dynamic programming: "overlapping-subproblem property"

Dynamic Time Warping



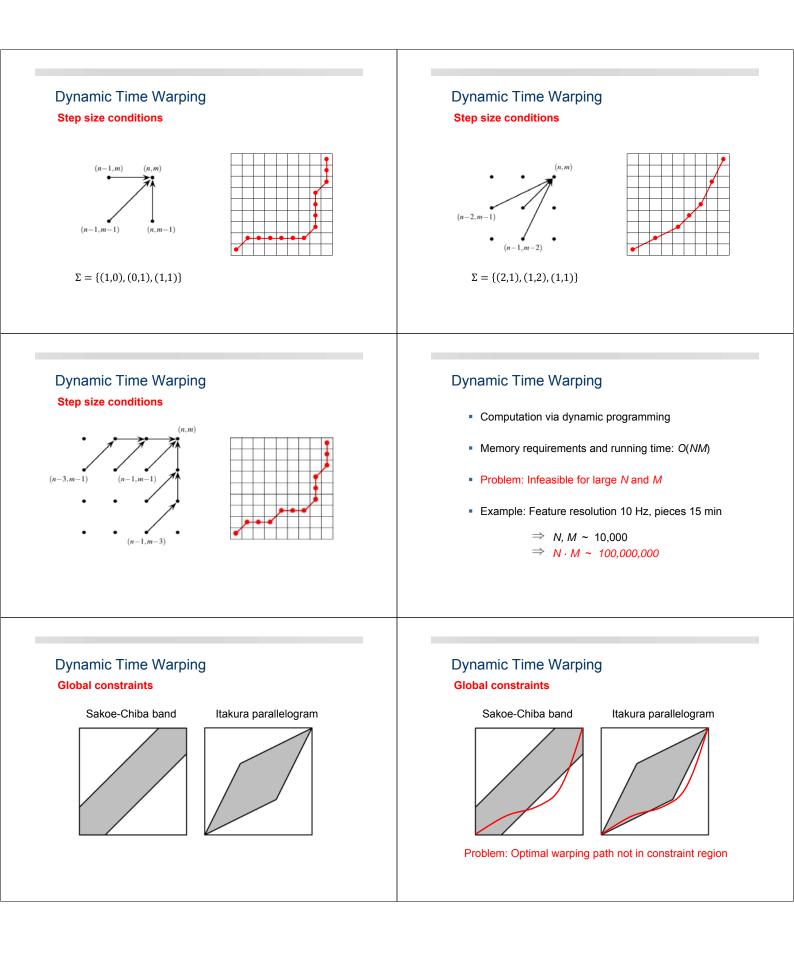
Dynamic Time Warping

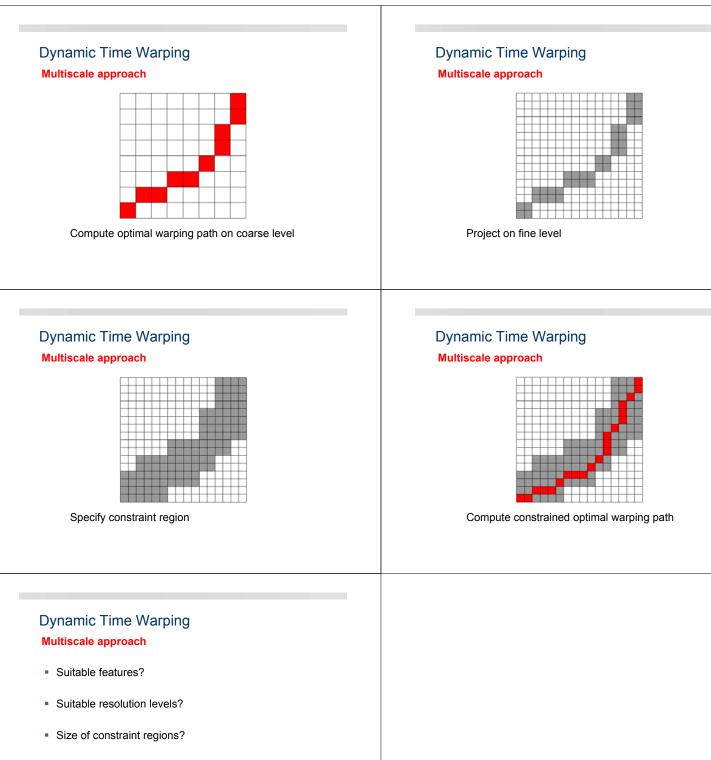
Example

X = (1,3,3,8,1) Y = (2,0,0,8,7,2) $c(x,y) = |x - y|, x,y \in \mathbb{R}$



Optimal warping path: $P^* = ((1,1), (2,2), (3,3), (4,4), (4,5), (5,6))$





Good trade-off between efficiency and robustness? Suitable parameters depend very much on application!